$$A = LDM^t$$

If A is nonsingular and A = LU, then we can set D = diag(U) and (since D is nonsinglar) $M^t \equiv D^{-1}U$ is a unit upper triangular matrix and $A = LDM^t$.

There is no inherent benefit to this factorization over LU, but it can give us a perspective from which to develop other algorithms. The idea is not to compute LDM^t from LU, but to derive a method to compute L, D and M directly. To that end, consider the k^{th} column of $A = LDM^t$:

$$a \equiv Ae_k = LDM^t e_k \equiv Ly. \tag{1}$$

Suppose we have already know the first k - 1 columns of L and consider the blocked treatment of the unit lower triangular system a = Ly:

$$\begin{bmatrix} L_{11} & 0\\ L_{21} & L_{22} \end{bmatrix} \begin{bmatrix} y_1\\ y_2 \end{bmatrix} = \begin{bmatrix} a_1\\ a_2 \end{bmatrix},$$
(2)

where L_{11} is $k \times k$ and *known*, but the last column of L_{21} is something we would like to compute. Forward substitution gives y_1 as the solution to $L_{11}y_1 = a_1$. Now we know the first k elements of y, and (from (1))

$$y = DM^t e_k, (3)$$

giving $e_k^t y = e_k^t D M^t e_k = d_{kk} e_k^t M^t e_k = d_{kk}$ (right?). $D^{-1}y = M^t e_k$ with M^t unit upper triangular means we can now compute the k^{th} column of M^t : $e_k^t M = [y_1^t D_1^{-1}, 0]$, where $D_1 = \text{diag}(d_{11}, d_{22}, \dots, d_{kk})$.

Now let's try to find z, the k^{th} column of $L_{21} = [\tilde{L}_{21}, z]$. Notice from (3) that M^t upper triangular means $y_2 = 0$, so $L_{21}y_1 + L_{22}y_2 = a_2$ from (2) reduces to $L_{21}y_1 = a_2$, or (with $y_1 = (\tilde{y}_1^t, d_{kk})^t$)

$$L_{21}\tilde{y_1} + zd_{kk} = a_2$$

giving $z = (a_2 - \tilde{L}_{21}\tilde{y}_1)/d_{kk}$.

Let's recap: Given the first k-1 columns of L and D, we can compute the k^{th} column of L, D and M^t as follows:

- 1. Solve $L_{11}y_1 = a_1$ for y_1 (forward substitution, about k^2 flops).
- 2. Set $d_{kk} = e_k^t y_1$.
- 3. Compute k^{th} row of M: $m_{kj} = e_j^t y_1/d_{jj}, \ j = 1, 2, ..., k-1 \ (k-1 \text{ flops}).$
- 4. Compute k^{th} column of L: $z = (a_2 \tilde{L}_{21}\tilde{y}_1)/d_{kk}$ (about 2k(n-k) flops).

Like the LU factorization, this takes about $2n^3/3$ flops. And like the LU factorization, the method is not a stable general purpose method (trouble if d_{kk} is |small|), but can be stabilized with little extra effort.

Unlike the LU factorization, this method can take advantage of symmetry. If $A = A^t$, then M = L and the $A = LDL^t$ factorization can be computed in $n^3/3$ flops, since step 4. can be skipped (it is done in step 3.).